

Lecture 7:

In discrete case, a differential equation can be discretized as:

where $\vec{u} = \begin{pmatrix} u(x_1) \\ u(x_2) \\ \vdots \\ u(x_N) \end{pmatrix}$ = values of u at N points $\{x_1, x_2, \dots, x_N\}$

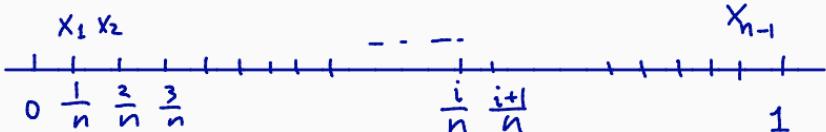
$\vec{g} = \begin{pmatrix} g(x_1) \\ g(x_2) \\ \vdots \\ g(x_N) \end{pmatrix}$ = values of g at N points $\{x_1, x_2, \dots, x_N\}$

D = $N \times N$ matrix approximating the differential operator.

Question: Can we "transform" \vec{u} and \vec{g} to turn the (BIG) linear system to SIMPLE algebraic equation?

Answer: YES! Discrete Fourier Transform !!

$$\frac{d^2 f}{dx^2}(x) = g(x) \quad 0 < x < 1 \quad \text{with } f(0) = 1 ; f(1) = 2 .$$



Discretize $(0, 1)$:

Approximation of $\frac{d^2 f}{dx^2}$:

$$f(x_{i+1}) \approx f(x_i) + \frac{1}{n} f'(x_i) + \frac{1}{2!} \frac{1}{n^2} f''(x_i) \quad (\text{Taylor's expansion})$$

$$+ f(x_{i-1}) \approx f(x_i) - \frac{1}{n} f'(x_i) + \frac{1}{2!} \frac{1}{n^2} f''(x_i)$$

$$f(x_{i+1}) + f(x_{i-1}) \approx 2f(x_i) + \frac{1}{n^2} f''(x_i)$$

$$\therefore \frac{d^2 f}{dx^2}(x_i) \approx \frac{f(x_{i+1}) - 2f(x_i) + f(x_{i-1})}{\left(\frac{1}{n^2}\right)}$$

$$\therefore \frac{d^2 f}{dx^2}(x_i) = g(x_i) \quad (\Rightarrow) \quad \left\{ \begin{array}{l} \frac{f(x_{i+1}) - 2f(x_i) + f(x_{i-1})}{\left(\frac{1}{n^2}\right)} = g(x_i) \\ \vdots \end{array} \right.$$

$$D \vec{f} = \vec{g} ; \vec{f} = \begin{pmatrix} f(x_1) \\ f(x_2) \\ \vdots \\ f(x_{n-1}) \end{pmatrix} ; \vec{g} = \begin{pmatrix} g(x_1) \\ g(x_2) \\ \vdots \\ g(x_{n-1}) \end{pmatrix} \quad \text{for } i=1, 2, \dots, n-1$$

Question: Extension to discrete case (Computational Math.)

Answer: Discrete Fourier Transform

- Goal:
- ① Define discrete Fourier Transform (DFT)
 - ② Use DFT to solve discretized differential eqt.

Definition: (Discrete Fourier Transform) Given $f_0, f_1, \dots, f_{n-1} \in \mathbb{C}$, then the discrete Fourier Transform (DFT) is defined as:

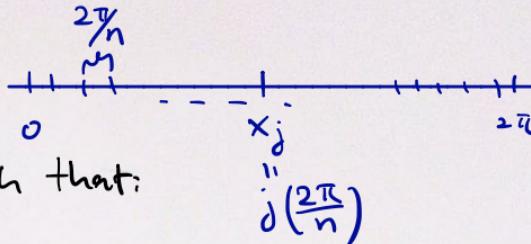
$$\vec{c} = \begin{pmatrix} c_0 \\ c_1 \\ \vdots \\ c_{n-1} \end{pmatrix} \in \mathbb{C}^n \text{ where } c_k = \frac{1}{n} \sum_{j=0}^{n-1} f_j e^{-i\left(\frac{2jk\pi}{n}\right)} \quad \text{for } k=0, 1, 2, \dots, n-1$$

The inverse discrete Fourier Transform recovers the original signal:

$$f_j = \sum_{k=0}^{n-1} c_k e^{i\left(\frac{2jk\pi}{n}\right)} \quad \text{for } j=0, 1, 2, \dots, n-1$$

Motivation 1: Let $f(x)$ defined on $[0, 2\pi]$

Approximate $f(x)$ by:



$$F_n(x) = \sum_{k=0}^{n-1} c_k e^{ikx}, \quad x \in [0, 2\pi] \text{ such that:}$$

$$F_n(x_j) = f(x_j) := f_j, \quad x_j = \frac{j\pi}{n}. \quad (\text{for all } j=0, 1, 2, \dots, n-1)$$

$$\left\{ \begin{array}{l} F_n(x_0) = c_0 + c_1 + c_2 + \dots + c_{n-1} = f_0 \\ F_n(x_1) = c_0 + c_1 e^{ix_1} + c_2 e^{i2x_1} + \dots + c_{n-1} e^{i(n-1)x_1} = f_1 \\ \vdots \\ F_n(x_{n-1}) = c_0 + c_1 e^{ix_{n-1}} + c_2 e^{i2x_{n-1}} + \dots + c_{n-1} e^{i(n-1)x_{n-1}} = f_{n-1} \end{array} \right.$$

Let $\omega = e^{i\frac{2\pi}{n}} = e^{ix_1}$, $\omega^2 = e^{i\frac{4\pi}{n}} = e^{ix_2}$, $\omega^3 = e^{ix_3}$, ... etc.

$\therefore (*)$ can be written as:

$$A_\omega \neq \begin{pmatrix} 1 & 1 & & & 1 \\ 1 & \omega & \omega^2 & \cdots & \omega^{n-1} \\ 1 & \omega^2 & \omega^4 & \cdots & \omega^{2(n-1)} \\ 1 & \vdots & & & \omega^{(n-1)^2} \\ 1 & \omega^{n-1} & \omega^{2(n-1)} & \cdots & \end{pmatrix} \begin{pmatrix} c_0 \\ c_1 \\ \vdots \\ c_{n-1} \end{pmatrix} = \begin{pmatrix} f_0 \\ f_1 \\ f_2 \\ \vdots \\ f_{n-1} \end{pmatrix}$$

$$\therefore (A_\omega \bar{A}_\omega)_{j,k} = 1 \cdot 1 + \omega^j \bar{\omega}^k + \omega^{2j} \bar{\omega}^{2k} + \dots + \omega^{(n-1)j} \bar{\omega}^{(n-1)k}$$

$$= 1 + e^{2\pi i \frac{(j-k)}{n}} + e^{\dots} + e^{2\pi i \frac{(n-1)(j-k)}{n}}$$

$$= \begin{cases} n & \text{if } j=k \\ \frac{1 - (e^{2\pi i \frac{(j-k)}{n}})^n}{1 - e^{2\pi i \frac{(j-k)}{n}}} = 0 & \text{if } j \neq k \end{cases}$$

$$\therefore \overline{A_w} \overline{A_w} = n I = \overline{A_w} A_w$$

We have:

$$\begin{pmatrix} c_0 \\ c_1 \\ \vdots \\ c_{n-1} \end{pmatrix} = A_w^{-1} \begin{pmatrix} f_0 \\ f_1 \\ \vdots \\ f_{n-1} \end{pmatrix} = \frac{\overline{A_w}}{n} \begin{pmatrix} f_0 \\ f_1 \\ \vdots \\ f_{n-1} \end{pmatrix}$$

$$\therefore C_k = \frac{1}{n} (f_0 + e^{-\frac{2\pi i}{n} k} f_1 + \dots + e^{-\frac{2\pi i}{n} k(n-1)} f_{n-1})$$

|| for $k=0, 1, 2, \dots, n-1$

Remark: Comp cost for DFT?

n^2 multiplication
 $(n-1)n$ addition $= O(n^2)$

Remark: Computational cost for DFT is:

$$\begin{array}{c} n^2 \text{ multiplication} \\ + \\ n(n-1) \text{ addition} \end{array} = \mathcal{O}(n^2)$$

Example: Consider $f(t) = 5 + 2\cos(t - \frac{\pi}{2}) + 3\cos(2t)$.

f is 2π -periodic. Divide $[0, 2\pi]$ by 4 partitions. Find the DFT of f (discretized by 4 points).

$$f_0 = f(0) = 8; f_1 = f\left(\frac{2\pi}{4}\right) = 4; f_2 = f\left(\frac{4\pi}{4}\right) = 8; f_3 = f\left(\frac{6\pi}{4}\right) = 0$$

$$\therefore \text{DFT: } C_k = \frac{1}{4} \sum_{j=0}^3 f_j e^{-i\left(\frac{2jk\pi}{4}\right)} \text{ for } k=0, 1, 2, 3 \quad \text{or}$$

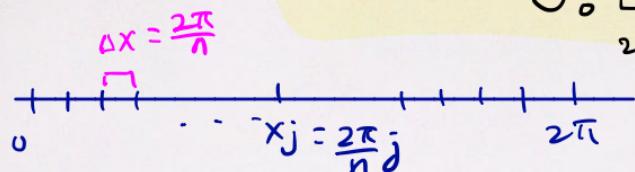
$$\begin{pmatrix} C_0 \\ C_1 \\ C_2 \\ C_3 \end{pmatrix} = \frac{1}{\sqrt{n}} \begin{pmatrix} f_0 \\ f_1 \\ f_2 \\ f_3 \end{pmatrix} = \frac{1}{\sqrt{4}} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -i & -1 & i \\ 1 & 1 & 1 & -1 \\ 1 & i & -1 & -i \end{pmatrix} \begin{pmatrix} 8 \\ 4 \\ 8 \\ 0 \end{pmatrix} = \begin{pmatrix} 5 \\ -i \\ 3 \\ i \end{pmatrix}$$

$\omega = e^{\frac{2\pi i}{4}}$

Motivation 2: Fourier Transform \leftrightarrow Fourier Series extended to related $(-\infty, \infty)$

Fourier coefficients = $C_k = \frac{1}{2\pi} \int_0^{2\pi} \underbrace{f(x)}_{2\pi\text{- periodic}} e^{-ikx} dx$

Divide:



We can approximate the integration:

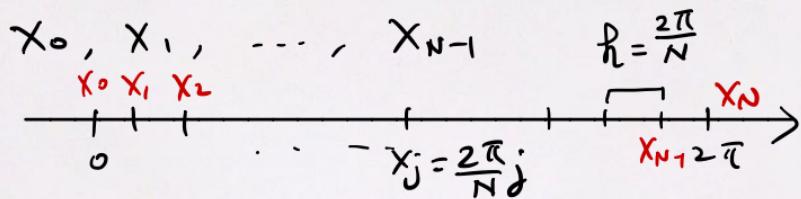
$$\begin{aligned} C_k &\approx \frac{1}{2\pi} \sum_{j=0}^{n-1} f(x_j) e^{-ikx_j} \Delta x = \frac{1}{2\pi} \sum_{j=0}^{n-1} f_j e^{-i\frac{2\pi}{n} j k} \left(\frac{2\pi}{n}\right) \\ &= \frac{1}{n} \sum_{j=0}^{n-1} f_j e^{-i\frac{2jk\pi}{n}} \quad \text{for } k=0, 1, 2, \dots, n-1 \\ &= DFT \end{aligned}$$

DFT = approximation of (complex) Fourier coefficient.

DFT and numerical diff eqt

Consider : $\frac{d^2 u}{dx^2} = f$ for $x \in [0, 2\pi]$ with periodic boundary condition $u(0) = u(2\pi)$

Suppose f is measured only at N discrete points :



Let $\vec{f} = \begin{pmatrix} f(x_0) \\ f(x_1) \\ \vdots \\ f(x_{N-1}) \end{pmatrix} \stackrel{\text{def}}{=} \begin{pmatrix} f_0 \\ f_1 \\ \vdots \\ f_{N-1} \end{pmatrix} \in \mathbb{R}^N$ and $\vec{u} = \begin{pmatrix} u(x_0) \\ u(x_1) \\ \vdots \\ u(x_{N-1}) \end{pmatrix} \stackrel{\text{def}}{=} \begin{pmatrix} u_0 \\ u_1 \\ \vdots \\ u_{N-1} \end{pmatrix} \in \mathbb{R}^N$

(unknown)

By Taylor's expansion,

$$u(x_j + h) \approx u(x_j) + hu'(x_j) + \frac{h^2}{2}u''(x_j) \quad (1)$$

$$u(x_j - h) \approx u(x_j) - hu'(x_j) + \frac{h^2}{2}u''(x_j) \quad (2)$$

$$(1) + (2) : u''(x_j) \approx \frac{u(x_j - h) - 2u(x_j) + u(x_j + h)}{h^2}$$

$$\therefore u''(x_j) \approx \frac{u_{j-1} - 2u_j + u_{j+1}}{h^2} \quad (\text{Central difference approximation})$$

Thus:

$$\begin{pmatrix} u''(x_0) \\ u''(x_1) \\ \vdots \\ u''(x_{N-1}) \end{pmatrix} \approx \tilde{D}\vec{u} \quad \text{where} \quad \tilde{D} = \frac{1}{h^2} \begin{pmatrix} -2 & 1 & & & & & & \\ 1 & -2 & 1 & & & & & \\ & \ddots & \ddots & \ddots & & & & \\ & & & & 1 & & & \\ & & & & & 1 & & \\ & & & & & & \ddots & \\ & & & & & & & -2 \end{pmatrix}$$

for $j = 0, 1, 2, \dots, N-1$

(Use the fact that $u_0 = u_N, u_{-1} = u_{N-1}$)

$\therefore \frac{d^2u}{dx^2} = f$ can be discretized as $\underset{M \times N}{\Rightarrow} \tilde{D} \vec{u} = \vec{f}$ (Linear System)

Numerical differential eqt

Remark: \tilde{D} is BlG matrix !!

Goal: Design numerical spectral method to solve $\tilde{D} \vec{u} = \vec{f}$.

Need to: Determine eigenvalues / eigenvectors of \tilde{D}